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Chemical Reaction Optimization for Population Transition in Peer-to-Peer Live Streaming

Albert Y.S. Lam, Member, IEEE, Jialing Xu, Student Member, IEEE, and Victor O.K. Li, Fellow, IEEE

Abstract—Peer-to-peer (P2P) live streaming applications are very popular in recent years and a Markov open queueing network model was developed to study the population dynamics in P2P live streaming. Based on the model, we deduce an optimization problem, called population transition problem, with the objective of maximizing the probability of universal streaming by manipulating population transition probability matrix. We employ a chemical reaction-inspired metaheuristic, Chemical Reaction Optimization (CRO), to solve the problem. Simulation results show that CRO outperforms many commonly used strategies for controlling population transition in many practical P2P live streaming systems. This work also shows that CRO also demonstrates the usability of CRO to solve optimization problems.

I. INTRODUCTION

Peer-to-peer (P2P) streaming has become one of the killer applications in the Internet. Many companies providing commercial P2P streaming services have been established, including UUSee [1], PPlive [2], SopCast [3], etc. For example, iResearch [4] shows that PPlive services have access rates of more than 10 million per month. The huge demand of P2P streaming services creates many research problems, including scheduling [5], topology control [6], and QoS support [7].

In addition to traditional video broadcasting, P2P technology provides a new platform for live streaming with less stringent geographical constraint and at lower costs. However, the stringent time requirement of P2P live streaming creates extra challenges. Xu and Li are the first to analyze the population dynamics for live streaming over P2P networks [8]. They propose a stochastic model and a new metric is introduced to evaluate the performance of different population transition strategies.

In this work, based on [8], we formulate the model as an optimization problem, entitled the Population Transition Problem (PTP). Solution of this problem requires exponentially increasing computation time with the problem scale. Due to its complexity, it can be considered as a black box and traditional optimization techniques are not effective. Therefore, we rely on evolutionary approaches. A newly proposed metaheuristic called Chemical Reaction Optimization (CRO) [9] is adopted to tackle PTP.

CRO is a chemical reaction-inspired general-purpose optimization method and it mimics the collisions of molecules in chemical reactions. Molecular interactions bring molecules from high to low energy states. This tendency allows us to explore the global optimum of the solution space. CRO has been shown to be effective in many applications [9], [10]. This motivates us to apply CRO to solve PTP. Simulation results show that CRO outperforms other commonly used strategies.

CRO is an evolutionary algorithm. Other evolutionary algorithms have been applied to P2P, e.g. [11] and [12], but they do not consider PTP. Instead, [11] focuses on constructing decentralized and self-organized P2P information system in computational grids and [12] addresses the packet scheduling issue for QoS support in P2P video streaming. Moreover, P2P has also been used as a platform to operate evolutionary algorithms, e.g. [13], [14], [15], but this is not the focus of this paper.

The rest of the paper is organized as follows. In Section II, we formulate the problem deduced from the model given in [8]. Section III describes the framework of CRO. We give the simulation results for performance evaluation in Section IV. Finally, we conclude the paper and suggest possible future work in Section V.

II. PROBLEM FORMULATION

In this section, we firstly revise the P2P streaming systems and the population dynamics model. Then we formulate the problem in the form of maximization.

A. Peer-to-Peer Live Streaming Systems

A single-channel P2P live streaming system consists of a stream source and a group of peers. The stream source provides real time streaming data (e.g. live video broadcast) to the peers. The peers may connect to networks with different topologies and traffic patterns, and thus, their cumulative transmission delays (total delay from the moment when the source sends out a packet to when the peer receives it) vary. As pointed out in [8], measuring the exact locations of peers are infeasible. So we differentiate the peers according to their cumulative transmission delays. Two peers with smaller and larger delays are called upstream and downstream peers, respectively. Upstream peers receive streaming data from the source earlier than downstream peers, act like a source and provide streaming data to its downstream peers. More upstream peers help support more downstream peers. Therefore, the total capacity of the system can improve. This demonstrates the power of P2P system compared with the traditional client-server architecture [16].

Many systems have tracker servers which provide system information to the peers. The most important information is a list of addresses of upstream peers for a peer to connect.
Fig. 1. Modeling colonies of peers as a Markov open queueing network

with. It is particularly useful when a peer joins a system. When a peer experiences degraded services, a tracker server may also provide it a list of recommended peers to re-attach to. Depending on the system implementation, this may be automatically completed by the streaming software. Thus, a system may control the topology of the P2P overlay network so as to improve the system performance.

B. Population Dynamics Model

We divide the peers into several colonies according to their cumulative transmission delays. Each peer joins a particular colony when participating in the streaming system and those with similar delays are put into the same colony. A peer may leave its colony and attach to another upon sensing its accumulative transmission delay has changed\(^1\). It may leave the system forever if it is not interested in the channel anymore.

The system is modeled as a Markov open queueing network (see Fig. 1). A peer is just like a customer in the network of queues. A peer participating in a colony, sojourning within a colony, shuttling between colonies, and quitting the system is tantamount to a customer entering a queue, getting served, transiting to other queues, and leaving the network, respectively. Here we consider a special case of the model with Poisson arrivals and exponential service times. We assume that the queues are \(M/M/\infty\) as a peer immediately gets served without waiting once it joins a colony. We also assume that there are \(n\) colonies and peers join the \(i\)-th colony from outside the system as a Poisson process with arrival rate \(v_i\) for \(i = 1, \ldots, n\). Peers stay at any colony for an exponentially distributed time period with mean \(\frac{1}{\mu_i}\). After a peer has been served, it may stay in or quit the system with probability \(P_{si}\) and \((1 - P_{si})\), respectively. If it stays, the probability of moving from the \(i\)-th to the \(j\)-th colony is \(P_{ij}\). By Theorem 2 of [8], the system has one and only one equilibrium distribution if there exist positive \(\alpha_i\)'s satisfying

\[
\alpha_i \mu = v_i + \mu \sum_{j=1}^{n} \alpha_j P_{ji}, \quad i = 1, \ldots, n, \tag{1}
\]

where \(\alpha_i\) represents the average number of peers in the \(i\)-th colony. Let \(x_i\) be the population size (i.e. number of peers) of the \(i\)-th colony. By Theorem 3 of [8], the probability of having a population size \(x_i\) for the \(i\)-th colony, \(\pi_i(x_i)\), are independent of each other in the equilibrium state, and given by

\[
\pi_i(x_i) = e^{-x_i \frac{\alpha_i}{x_i^1}}, \quad i = 1, \ldots, n. \tag{2}
\]

Thus, the probability of the whole network in equilibrium state \(s = (x_1, \ldots, x_n)\) is given by

\[
\pi(s) = \prod_{i=1}^{n} \pi_i(x_i). \tag{3}
\]

A P2P streaming system is considered to be in universal streaming when all peers in the system can receive sufficient streaming data. To make the universal streaming condition tractable, we suppose that peers in the \(i\)-th colony can only get data supply from those in the same colony and their immediate upstream colony, i.e. the \((i - 1)\)-th colony. Moreover, the maximum number of peers with sufficient data supply in the \(i\)-th colony is determined by the population size in the \((i - 1)\)-th colony multiplied by a sufficient provision ratio \(b^1\), i.e.,

\[
x_i \leq \lfloor x_{i-1} \cdot b \rfloor.
\]

Let \(B\) be the capacity of the stream source. The probability of universal streaming is defined as

\[
P_{us} = \Pr\{x_1 \leq B, x_2 \leq \lfloor x_1 \cdot b \rfloor, \ldots, x_n \leq \lfloor x_{n-1} \cdot b \rfloor\}.
\]

With the consideration of system state probability given by (3), \(P_{us}\) can be further expressed as

\[
P_{us} = \sum_{s \in \Phi} \pi(s), \tag{4}
\]

where \(\Phi\) is the set of all system states satisfying universal streaming, i.e.,

\[
\Phi = \{(x_1, \ldots, x_n) | x_1 \leq B, x_2 \leq \lfloor x_1 \cdot b \rfloor, \ldots, x_n \leq \lfloor x_{n-1} \cdot b \rfloor\}.
\]

Here we only give an overview of the population dynamics model. For more information about the model, interested readers may refer to [8].

\(^1\)A peer which participates in an inappropriate colony may degrade the overall system performance. For example, a peer with high delay is not suitable to act as a source for transmitting data to peers with lower delay.

\(^2\)\(x_i\) must be a non-negative integer.

\(^3\)\(b\) usually has a value between one and two in practice.
\[
\begin{bmatrix}
v_1 \\
v_2 \\
\vdots \\
v_n
\end{bmatrix} = \mu
\begin{bmatrix}
(1 - P_s P_{11}) & -P_s P_{21} & \cdots & -P_s P_{n1} \\
-P_s P_{12} & (1 - P_s P_{22}) & \cdots & -P_s P_{n2} \\
\vdots & \vdots & \ddots & \vdots \\
-P_s P_{1n} & -P_s P_{2n} & \cdots & (1 - P_s P_{nn})
\end{bmatrix}
\begin{bmatrix}
\alpha_1 \\
\alpha_2 \\
\vdots \\
\alpha_n
\end{bmatrix}
\]

\eqref{eq:6}

\section*{C. Population Transition Problem}

The system parameters \( B, b, n, \mu, \) and \( v_i, i = 1, \ldots, n, \) characterize the P2P streaming system. Recall that the system can control the topology of the overlay P2P network for providing better service. When a peer needs to switch from a colony to another, the system can determine to which colony it should connect (by controlling the list of available upstream peers provided to the peer). One simple implementation is to control the population transition strategy, i.e. the population transition probability matrix \([P_{ij}]\) \(1 \leq i, j \leq n\) (or \([P_{ij}]\) in short), so as to maximize the probability of universal streaming \(P_{us}\).

Here we try to demonstrate how to compute \(P_{us}\) corresponding to a matrix \([P_{ij}]\). Recall that \(P_{ij}\) is the probability of a peer transiting from the \(i\)-th to the \(j\)-th colony. Therefore, we have

\[
\sum_{j=1}^{n} P_{ij} = 1, \quad i = 1, \ldots, n, \quad \text{(5)}
\]

and it is a constraint of PTP. From \(1\), we have

\[
v_i = \mu (\alpha_i - \sum_{j=1}^{n} \alpha_j P_s P_{ji}), \quad i = 1, \ldots, n,
\]

which in matrix form gives \(6\) or

\[
v = \mu (I - P_s P^T) \alpha, \quad \text{(7)}
\]

where \(v\) is the arrival rate vector \([v_1, \ldots, v_n]^T\), \(I\) is an identity matrix of size \(n\), \(P\) is the transition matrix of the Markov chain \([P_{ij}]\), and \(\alpha\) equals \([\alpha_1, \ldots, \alpha_n]^T\). We need to solve the system of linear equations \(7\) to get \(\alpha\). Let \(M\) and \(\tilde{M}\) be the matrix \(\mu (I - P_s P^T)\) and the augmented matrix of \(M\), respectively, i.e.,

\[
\tilde{M} =
\begin{bmatrix}
\mu(1 - P_s P_{11}) & -\mu P_s P_{21} & \cdots & -\mu P_s P_{n1} \\
-\mu P_s P_{12} & \mu(1 - P_s P_{22}) & \cdots & -\mu P_s P_{n2} \\
\vdots & \vdots & \ddots & \vdots \\
-\mu P_s P_{1n} & -\mu P_s P_{2n} & \cdots & \mu(1 - P_s P_{nn})
\end{bmatrix}
\begin{bmatrix}
v_1 \\
v_2 \\
\vdots \\
v_n
\end{bmatrix}
\]

By the Kronecker-Capelli Theorem \(17\), \(7\) has a solution for \(\alpha\) provided that

\[
\text{rank}(M) = \text{rank}(\tilde{M}). \quad \text{(8)}
\]

Therefore, \(8\) is also a constraint of PTP. We use \(7\) to compute \(\alpha\) from \(P\).

For a particular state \(s = (x_1, \ldots, x_n)\), we calculate its probability with \(2\) and \(3\). We determine \(P_{us}\) by summing the probabilities of all states satisfying universal streaming with \(4\). Consider an example when \(n = 2\), \(B = 5\), and \(b = 1.5\). The feasible states \((x_1, x_2)\) include:

\[
(5, 0), (4, 0), (3, 0), (2, 0), (1, 0), (0, 0), (5, 1), (4, 1), (3, 1), (2, 1), (1, 1), (5, 2), (4, 2), (3, 2), (2, 2), (5, 3), (4, 3), (3, 3), (2, 3), (5, 4), (4, 4), (3, 4), (5, 5), (4, 5), (5, 6), (4, 6), (5, 7).
\]

By inspection, the size of the feasible state space \(\Phi\) is expressed as

\[
\sum_{i=0}^{5} \sum_{j=0}^{5} \cdots \sum_{n-1=0}^{5} (1 + [i_{n-1}b]), \quad \text{(9)}
\]

where \([i_0b]\) is equal to \(B\). This shows that the computation of \(P_{us}\) grows exponentially with \(n\), \(B\), and \(b\).

A possible solution \(\omega\) of the problem is a probability transition matrix \([P_{ij}]\) and the corresponding objective function value \(y\) is its \(P_{us}\). The above process, denoted by \(f\), shows how to compute \(P_{us}\) for a particular \([P_{ij}]\), i.e.,

\[
P_{us} = f([P_{ij}]). \quad \text{(10)}
\]

It is the objective function which we need to maximize. \(9\) reveals that \(f\) is very computationally intensive when the scale of the P2P streaming system grows. To summarize, PTP is mathematically represented by

\[
\text{max } P_{us} = f([P_{ij}]) \quad \text{subject to } \sum_{j=1}^{n} P_{ij} = 1, \quad i = 1, \ldots, n, \quad \text{rank}(M) = \text{rank}(\tilde{M}).
\]

Different system characteristics (e.g. values of \(\mu\) and \(v\)) imply different problem instances of PTP, and thus, the optimal solutions corresponding to different instances may vary. In other words, we need to solve the PTP whenever the system parameters change. According to \(18\), the system characteristics follow similar daily patterns. Therefore, we can construct several instances of PTP with representative system parameter settings corresponding to such patterns, and the computed \([P_{ij}]\)’s can be utilized to optimize the system performance. We will present an algorithm to solve PTP in the next section.
III. ALGORITHM DESIGN

There are not many optimization problems which have right stochastic matrices as solutions. Thus, we introduce several operators to deal with the solutions satisfying the constraints posed on PTP. Then we discuss how CRO works.

A. Operators

1) Initial Solution Generator: This operator is used to generate the initial solutions for CRO to manipulate. Each call to the operator will generate a right stochastic matrix of order n. It is done by firstly generating a matrix of order n with element $P_{ij} \in [0, 1]$ randomly. Then we transform the elements so that the sum of each row is equal to one. For example,

$$\begin{bmatrix}
0.79 & 0.04 & 0.68 \\
0.96 & 0.85 & 0.76 \\
0.66 & 0.93 & 0.74
\end{bmatrix} \Rightarrow \begin{bmatrix}
0.53 & 0.02 & 0.45 \\
0.37 & 0.33 & 0.30 \\
0.28 & 0.40 & 0.32
\end{bmatrix}$$

Its pseudocode is given in Fig. 2.

2) Neighborhood Solution Generator: This operator is used to generate a new solution $\omega'$ in the neighborhood of a given solution $\omega$. It is done by randomly choosing a row $i$ ($1 \leq i \leq n$) of $\omega$ first. Then we select two distinct elements $P_{ij}$ and $P_{ik}$ ($1 \leq j, k \leq n$) in row $i$. Next we divide the sum of $P_{ij}$ and $P_{ik}$ into two random portions and assign each of them to $P_{ij}'$ and $P_{ik}'$ of $\omega'$, respectively. The rest of the elements in $\omega'$ are just the same as those in $\omega$. For example, $P_{12}$ and $P_{13}$ are chosen in the following:

$$\begin{bmatrix}
0.53 & 0.02 & 0.45 \\
0.37 & 0.33 & 0.30 \\
0.28 & 0.40 & 0.32
\end{bmatrix} \Rightarrow \begin{bmatrix}
0.53 & 0.16 & 0.31 \\
0.37 & 0.33 & 0.30 \\
0.28 & 0.40 & 0.32
\end{bmatrix}$$

Its pseudocode is given in Fig. 3.

3) Decomposition: This operator is used to produce two new solution $\omega_1'$ and $\omega_2'$ from a given solution $\omega$. For each row of $\omega$, we assign it to the same row of either $\omega_1'$ or $\omega_2'$ randomly. For those rows in $\omega_1'$ or $\omega_2'$ which have not been assigned with values, we make them rows of random positive values whose total is equal to one. This mechanism is similar to that used in the initial solution generator to produce a random row of numbers. For example,

$$\begin{bmatrix}
0.53 & 0.02 & 0.45 \\
0.37 & 0.33 & 0.30 \\
0.28 & 0.40 & 0.32
\end{bmatrix} \Rightarrow \begin{bmatrix}
0.53 & 0.02 & 0.45 \\
0.34 & 0.62 & 0.04 \\
0.28 & 0.40 & 0.32
\end{bmatrix} \text{ AND } \begin{bmatrix}
0.84 & 0.04 & 0.12 \\
0.37 & 0.33 & 0.30 \\
0.13 & 0.22 & 0.65
\end{bmatrix}$$

Its pseudocode is given in Fig. 4.

4) Synthesis: This operator is used to produce a new solution $\omega'$ from two solutions $\omega_1$ and $\omega_2$. For each row of $\omega'$, we assign it a corresponding row from either $\omega_1$ or
the second kind is CollMole. We define four types of elementary reactions, i.e. on-wall ineffective collision, decomposition, intermolecular collision, and synthesis, to characterize different collisions. The former two can be triggered when a molecule hits a wall. The latter two can take place when molecules collide with each other (we assume only two molecules are involved in the simulation). We follow the framework described in [9] to implement the algorithm. Interested readers can refer to [9] to comprehend how the whole algorithm works. To make it suitable for solving PTP, we have several modifications to the operators used in CRO. We adopt the initial solution generator described in Section III-A.1 to produce PopSize number of molecules at initialization. Their initial KEs are assigned a value equal to InitialKE. The neighborhood solution generator given in Section III-A.2 is utilized in on-wall and inter-molecular ineffective collisions to generate new solutions in the neighborhoods of existing ones. We apply the operators from Section III-A.3 and Section III-A.4 to perform decomposition and synthesis, respectively.

IV. SIMULATION

In this section, we firstly explain how we handle the constraints and then compare our solution with a set of common benchmarks. Next we describe the configuration of the simulation and finally discuss the results.

A. Constraints Handling

We have two types of constraints, i.e. (5) and (8). We try to confine the search to the feasible solution space. In other words, we keep generating solutions without violating any constraints in both initialization and iterations of CRO. To do this, we handle the two types of constraints separately. For (5), the initial solution generator ensures that all solutions it produces are right stochastic matrices. Moreover, the neighborhood solution generator, decomposition and synthesis operators are specially designed to transform right stochastic matrices to right stochastic matrices as well. For (8), with the possible values (non-negative real numbers) assigned to \( v, \mu, P_s, \) and \( P \), it is uncommon to generate a \([P_{ij}]\) which makes the system of linear equations (7) inconsistent\(^3\). If it does happen, we can simply ignore that \([P_{ij}]\) and generate a new one instead.

B. Benchmarks Problems

To have fair comparisons of performance over various strategies, we need a set of common benchmark problems in the simulation. Let \( V_T \) be the overall external arrival rate to the system, i.e. \( V_T = v_1 + v_2, \ldots, v_n \). We randomly generate 60 problem instances with \( n = 4 \) and \( V_T = 2, 4, \ldots, 20 \). For each \( V_T \), there are 6 problem instances, each of which is generated by randomly dividing \( V_T \) into four parts, representing \( v_1, v_2, v_3, \) and \( v_4 \), respectively. Each instance is represented by a number formatted in the form

\(^3\)We seldom encounter \([P_{ij}]\) making (7) inconsistent in the mass simulation.

---

\[
\begin{align*}
\text{procedure } & \text{Synthesize}(\omega_1, \omega_2) \\
& \text{for } i := 1 \text{ to } n \\
& \quad \text{Generate a random number } t \in [0, 1] \\
& \quad \text{if } t < 0.5 \\
& \quad \quad [P_{11}^n, \ldots, P_{in}^n] := [P_{(1)1}, \ldots, P_{(1)n}] \\
& \quad \text{else} \\
& \quad \quad [P_{11}^n, \ldots, P_{in}^n] := [P_{(2)1}, \ldots, P_{(2)n}] \\
& \quad \text{end if} \\
& \text{return } \omega' \\
& \text{end procedure}
\end{align*}
\]

Fig. 5. Pseudocode of synthesis

\[\begin{bmatrix}
0.53 & 0.02 & 0.45 \\
0.37 & 0.33 & 0.30 \\
0.28 & 0.40 & 0.32
\end{bmatrix}
\begin{bmatrix}
0.34 & 0.62 & 0.04 \\
0.84 & 0.04 & 0.12 \\
0.13 & 0.22 & 0.65
\end{bmatrix}
\overset{\text{AND}}{\Rightarrow}
\begin{bmatrix}
0.34 & 0.62 & 0.04 \\
0.84 & 0.04 & 0.12 \\
0.28 & 0.40 & 0.32
\end{bmatrix}
\]

\(\omega_2\) randomly. For example,
of “V_{T_case number”}. For example, “20.5” means that it is the fifth instance with V_{T equal to 20.

C. Implementation Details

The simulation codes are programmed in MATLAB because MATLAB has powerful and efficient operators to handle matrices and to solve (7) [19]. In computing (10), we follow [8] to set the system parameters, i.e. \( B = 15, \ b = 1.5, \ P_s = 0.6, \) and \( \mu = 0.5. \) For all the dynamic strategies (introduced later), the stopping criterion is when the maximum number of function evaluations (FEs), set to 1000, is reached. After several trial runs of the simulation, the parameter values of CRO are set as follows: \( \text{PopSize} = 15; \) \( \text{KELossRate} = 0.2; \) \( \text{MoleColl} = 0.3; \) \( \text{InitialKE} = 10; \) \( \alpha = 100; \) and \( \beta = 0.5. \)

PTP is a maximization problem. However, CRO is originally designed to solve minimization problems and the objective function value is interpreted as energy, which is of non-negative value. The typical way of converting a maximization problem \( f \) to a minimization one (i.e. letting \( -f \) as the objective function) may not be appropriate. Instead, we can consider \( f' = \text{offset} - f, \) where \( \text{offset} \) is supposed to be large enough to make every possible \( f' \) non-negative. After minimizing \( f' \), we can compute the corresponding \( f \) by \( f = \text{offset} - f' \). In this simulation, as \( P_{us} \) is upper bounded by one, we set the \( \text{offset} \) equal to one accordingly. This also demonstrates a technique on how to employ CRO to solve maximization problems.

D. Results

We compare the performance of CRO with other strategies which also manipulate \( [P_{ij}] \) as a solution, i.e. random search (Random), inert staying (IS), shorter delay (SD), and random walk (RW). CRO and Random are considered as dynamic approaches as they produce \( [P_{ij}] \) with respect to the problems, i.e. different \( [P_{ij}] \) are produced according to the values of \( \{v_1, \ldots, v_n\}. \) Random is introduced to serve as a reference for the dynamic approaches. IS, SD, and RW (introduced in [8]) are regarded as static strategies since they always use the same \( [P_{ij}] \) for different problems. Despite this, they still have practical use [8], [20], [21]. All simulations are performed on the same PC with Intel Core 2 Quad Processor Q9650 and 3GB of RAM. Each run of CRO and Random takes approximately 8.45s and 8.20s, respectively, and thus CRO suffers a little increase in computation time. As IS, SD, and RW always retain the same solution, one FE is required and they consume around \( \frac{8.20}{1000} \) s of computation time in each simulation run.

The simulation results for all problem instances are shown in Fig. 6. Each data point of CRO and Random is the mean of the results of 50 runs\(^6\) while the plus and minus error bars indicate the best and worst results of all 50 runs, respectively. The upper and lower rows of numbers in each chart show the standard deviations of the 50 runs for CRO and Random, separately. For some instances (e.g. 8.6, 10.6, 12.1, etc.), CRO performs almost the same as Random because these instances are particularly hard to solve, and thus, CRO does not show superiority over Random on the average. We can see that, in general, dynamic approaches can produce better results than the static ones. Moreover, CRO is superior to the others in most cases with less fluctuations across the instances of the same \( V_T. \) This reveals that CRO is more reliable to generate good solutions when solving new problem instances.

We also study the average percentage (%) improvement of CRO over other strategies, which is computed by

\[
\frac{\text{Result}_{CRO} - \text{Result}_{\xi}}{\text{Result}_{\xi}} \times 100\%,
\]

where \( \xi \) is any algorithm other than CRO. In Table I, each value is computed with the mean for all instances with respect to the same \( V_T. \) Although the “more streaming supply” (MSS) strategy [8] does not manipulate \( [P_{ij}] \) and thus it does not fit into the framework of PTP, we also include it in the table for completeness, as MSS is similar to the coolstreaming scheme [22]. CRO is obviously much better than IS and RW and it also outperforms the rest on the average.

For the dynamic approaches, we also give their convergence curves in Fig. 7, which shows the traces of best solutions found in a particular run. We only show the curves for the cases where CRO has best (20.5), median (2.4), and worst (12.1) performance over Random. We record one data point every 50 FEs. At the beginning (from 1 to 100 FEs, CRO may be inferior to Random because the performance at this stage highly depends on the initial solution. However, CRO gradually improves the solutions and get better ones after 100 FEs. Fig. 7 also shows that the stopping criterion used in the simulation is appropriate.

<table>
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<tr>
<th>( V_T )</th>
<th>Random</th>
<th>IS</th>
<th>SD</th>
<th>RW</th>
<th>MSS</th>
</tr>
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<tbody>
<tr>
<td>2</td>
<td>37.03</td>
<td>234.82</td>
<td>120.65</td>
<td>131.69</td>
<td>1.61</td>
</tr>
<tr>
<td>4</td>
<td>31.21</td>
<td>640.37</td>
<td>175.51</td>
<td>197.60</td>
<td>13.84</td>
</tr>
<tr>
<td>6</td>
<td>15.14</td>
<td>1421.87</td>
<td>175.82</td>
<td>186.66</td>
<td>24.28</td>
</tr>
<tr>
<td>8</td>
<td>6.29</td>
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\(^6\)CRO and Random are stochastic methods and the result produced in different runs may vary.

V. Conclusion

P2P live streaming is one of the killer applications in the Internet and the increasing size of P2P live streaming systems creates many challenging problems. Thus a Markov open queueing network model was developed to study the
Fig. 6. Simulation results
We will also try to apply CRO to solve other problems in annealing, particle swarm optimization, etc. in solving evolutionary algorithms (e.g. genetic algorithm, simulated annealing, particle swarm optimization, etc.) in solving optimization problems. In the future, we will demonstrate the capability of CRO in P2P live streaming and successfully applying CRO to solve the problem. This work also demonstrates the capability of CRO in solving optimization problems. Our contribution includes: 1) proposing PTP for P2P live streaming systems.

P2P systems.

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